

# A Comprehensive Research Study on Low-Interaction Secure Shell Honeypot

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# Abstract

This paper details information acquired from a secure shell honeypot, including plaintext login credentials and comprehensive attack data. As the number of data breaches and password leaks rises year after year, more dictionaries of reverse-engineered hashed passwords develop. Besides contributing to educational password this article also attempts to provide dictionaries, information about the geographical makeup of hackers encountered, as well as favored protocols. Its goal is to encourage developers to produce practical honeypot solutions for organizations with limited resources for cyber-protection, well their as as to encourage organizations to implement such measures and study their data. The low-interaction, user-friendly honeypot created capable of running without manual is intervention, and without interfering with parallelly running processes. Besides collecting login credentials used with SSH, in plaintext, its capabilities include recording, analyzing, and sending notifications about suspicious network traffic.

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### 1. Introduction

#### 1.1. Background

A network is a set of devices that use communication protocols to share resources. It establishes an architecture that allows a variety of equipment types to organize, unify and control hardware and software components of the network.

While networks have brought humanity closer than ever, their improper implementation or inadequate security can have very serious real-world consequences [1, 2], such as the remote deployment of computer viruses and worms, or the launch of Denial of Service (DoS) attacks.

Network security refers to the protection of data from unauthorized access, damage and development, and the implementation of policies and procedures for recovery from breaches and data losses. It can be implemented via an offensive approach, a defensive approach, or a hybrid approach. While offensive security is realised by deploying a proactive approach to security through the use of ethical hacking, defensive security uses a reactive approach to security that focuses on prevention, detection, and response to attacks.

Honeypots are emulated deceptive systems that can be used to assess where hackers infiltrating a network are coming from, the level of threat, their modus operandi, data of interest and the effectiveness of the hosting party's security stack. They are designed to trick the attacker into thinking a genuine system has been pawned, by purposely engaging them and identifying malicious activities performed by them over the internet. Honeypots are deliberately configured with known vulnerabilities in place, to make attractive targets for attackers. Since no interaction with a honeypot is authorized, all traffic is suspicious. Honeypots can thus automatically and accurately detect, analyze, and defend against zero-day and advanced attacks - providing insight into malicious activity within networks using a preventive, deceptive approach to security. The usage of tactics that rely on a thorough understanding of the system environment and its analysis to detect potential flaws influences the development and 100

deployment of preventive and protective measures that discourage or eliminate cyberattacks to a large extent. Due to this reason, honeypots are now being used in both, governmental and nongovernmental organisations such as banks, industrial control systems, educational institutions, etc.

### 1.2. Related Work

As defined by Joshi and Sardana [3], a honeypot is "A program that takes the appearance of an attractive service, set of services, an entire operating system or even an entire network, but is in reality, a tightly sealed compartment built to lure and contain an attacker". Covered by Tsikerdekis et al [4], most of the work available today concentrates on the development of unique honeypots that frequently target a specific feature, without offering a comprehensive understanding of how they might be built to prevent detection by attackers.

As summarised by Campbell et al [5], honeypots can be classified as (i) low-interaction, medium-interaction or high-interaction, on the basis of their functionality and supported services, (ii) deception, intimidation or reconnaissance on the basis of their mode of deployment, or (iii) production and research, on the basis of their deployment category. By conducting a comprehensive analysis of existing honeypot literature, they concluded that by the early 21st century, developed countries such as the United States of America and South Africa had provided far more insights into the usage and significance of honeypots than other countries, possibly due to their higher level of dependence on computing networks for daily functioning in those times.

Their insights made it evident that most of the research in this field took place when (i) internet usage started to grow in the absence of security standards (2002-2003), and (ii) internet-supported devices became commonplace, which led to its utilization for a diverse range of activities such as business, banking, social networking and the like (2006-2012). Themes such as new types of honeypots, improving the accuracy in threat detection, lowering false positives and avoiding detection appeared to be preferred over studies on the ethics of honeypots, mainly by researchers motivated by academic incentives that come with journal publication. Further explained by Tsikerdekis et al [4] and summarized in Table 1, honeypots that follow the Secure Shell (SSH) protocol without allowing much shell functionality and allow interactions for limited periods of time can be classified as low-interaction honeypots, usually placed in networks not being monitored by Intrusion Detection Systems (IDS). They are prone to detection and are configured as such. High-interaction honeypots, however, are configured to avoid detection to discover zero-day attacks and the modus operandi of hackers. For this reason, they emulate systems very thoroughly. This functionality legitimate is determined by the deployment category, i.e., research or production. While the former is placed within the network's Demilitarized Zone [6] to gather a wide range of threat intelligence, the latter maintains proximity to real assets for very specific intelligence from both, internal and external threats.

Anti-detection mechanism	Characteristic			
	Type (Research/ Production)	Interaction (Low/ High)	Scalability (Low cost/ High cost)	Implementation (Software/Hard ware)
Automatic honeypot redeployment	Either	Low	Low cost	Software
Honeypot delay reduction	Either	Either	Low cost	Software
Honeypot process transparency	Research	High	Low cost	Software
Dedicated hardware	Production	Either	High cost	Hardware

Dynamic	Research	High	High cost	Software
intelligence on honeypots				
<i>v</i> 1				

Table 1. Analysis of related works

Depending on the type of implementation, i.e. (i) hardware regular computers or specialized Supervisory Control and Data Acquisition (SCADA) systems, (ii) software-simulated hardware using virtualization, or (iii) hybrid, the scalability of honeypots becomes a notable factor, especially in the case of botnets, and/or state-sponsored attacks.

Exploring the theme of avoiding honeypot detection, this study laid out possible approaches that can be studied and implemented for more realistic emulations. Proposing (i) automatic honeypot redeployment - redeployment of the honeypot with an altered configuration upon detection by an attacker, (ii) honeypot delay reduction - minimization of delays caused by event logging - prone to detection unless the latency of the virtual honeypot network is lowered to match a physical network's link latency, (iii) honeypot process transparency - hiding unrealistic modified sequences of events such as the forwarding of connections between a honeypot's frontend and backend, by emulating a three-way Transmission Control Protocol (TCP) handshake while hiding the same, (iv) dedicated hardware - using specific hardware components to reduce software delays, increase system security, and enabling the system to support honeynets; and (v) dynamic intelligence on honeypots -the usage of machine learning and artificial intelligence to disable unexpected programs, dynamically change directory structures to increase attractiveness, and encourage attackers to reveal their geo-cultural identities on the basis of their interactions; the authors concluded that while a honeypot environment's alignment with an attacker's expectation of legitimate systems determines the chances of detection, constraints such as available hardware, development and maintenance costs, and legal restraints don't enable developers to build extremely efficient honeypots.

While these studies explore past literature and future implementation strategies in detail, the challenge of minimizing detection also depends on a thorough understanding of the challenges that require these solutions in the first place. Prior to the study by Tsikerdekis et al [4], Du [7] conducted research on the same, determining that honeypots mainly face issues in (i) hiding capture tools while collecting as much data as possible, (ii) capturing session data encrypted on the hacker's side, and (iii) collecting and transmitting data through secret channels. To combat the same, they proposed the following solutions: (i) Capture Tool Hiding via a) Module Hiding - deleting the pointer of the capture module of any data capturing tool loaded to the Linux kernel upon system initialization, and b) Process Hiding - changing the system call used to query process information in a system using the "ps" command, in order to stop programs using the system call from accessing the file, thus hiding the process. This can be effective as the program(s) within the honeypot would be executing multiple system processes; (ii) Session Encryption Data Capture while the execution of Trojan shells upon logging in can be exposed easily, changing the index of pointers of system calls such as read() and write() can enable the implementation of the capture module's own functions, which would result in direct access to the data that is part of such system calls, and (iii) Establishment of Hidden Data Transmission Channel - hiding the transfer of logs to centralized honeypot servers by configuring the capture module to transfer data via User Datagram Protocol (UDP) streams after altering the kernel on each endpoint such that data packets cannot be accessed. This would require the capture module to match the preset destination UDP port and magic number (a constant numerical value used to identify different protocols) on the endpoints within the Local Area Network (LAN) in order to make network sniffers on the endpoints ignore the packets.

Although this study was highly specific and dealt with issues directly at the kernel level, the highlighted approaches have certain drawbacks: (i) the capture module cannot be unloaded once it has been loaded, and the root user cannot locate it, and (ii) if the capture module contains a bug, the kernel may become unstable and the system may crash. These issues may have an impact on the normal operation of the honeypot, as well as the overall performance of the honeynet. The lack of implementation of these

suggestions provides no insight into the feasibility of these methods, especially in the long term.

Finally, recent comprehensive surveys [8,9, 10, 11] of the research on honeypots and honeynets for Internet of Things (IoT), Industrial Internet of Things (IIoT), and Cyber-Physical Systems (CPS) over the period 2002-2020 dealt with the taxonomy and analysis, key design factors, and open issues for future honeypots and honeynets for IoT, IIoT, and CPS environments revealed that the key to the design and implementation of competent honeypots lies in a good understanding of its target application area, purpose, cost, deployment location, intended level of interaction with the attacker, resource level, services, simulation or emulation, realistic service to the attacker, tools that will be used, the possibility of fingerprinting and indexing, and the liability issues that may come up.

To conclude, attackers have been able to detect honeypots and identify ways to exploit them because of

- the lack of research and expertise in emerging domains such as machine learning, unexplored protocols, anti-detection mechanisms, optimized deployment location, and the constant threat of insider attacks, and hardware vulnerabilities
- to date, much of the research has been focused on the creation of unique honeypots that typically focus on a single component without offering a comprehensive knowledge of how they could be structured to prevent detection by attackers
- the data been collected with certain restrictions, such as short time ranges, cultural biases, a narrow range of tools/technologies tested, etc.
- the large majority of these honeypots are built on outdated systems, with poor maintenance and irregular development cycles. Accessible to both, security professionals and attackers, they are predictable due to their limited adaptability and poor deception [4].

The integration and expansion of these categories could provide a clearer understanding of current issues, and the methods of eradicating them.

Proposed solutions are either valid under very strict conditions - on the basis of necessary hardware and software - or aren't comprehensive of the above-mentioned factors. Additionally, for a honeypot to be feasible and effective, a certain degree of deception is absolutely necessary, which isn't provided by the default configurations of most non-commercial honeypots.

# 2. Problem Statement

As mentioned earlier, the primary limitation of currently available honeypots lies in their deception capabilities, and the level of technical knowledge required for their efficient usage. In today's highly connected and extremely vulnerable digital space, honeypots are a necessary defence mechanism not only for niche research institutions and/or large organisations with a considerable security-focused workforce but also for smaller organisations dependent on the internet for any degree of daily functioning - regardless of their technical expertise [9]. Thus, arises the problem statement, and the proposed solution:

"The availability of open-source honeypots makes defensive network security easier for organisations across industries. However, the level of technical expertise required to customise their configuration and improve their deception abilities is not available to small organisations. This gap in requirement vs availability means that the advancement in honeypot research has not yet resulted in enough real-world implementation of proposed deception solutions to make this technology feasible for the global community. To minimise the need for small organisations to have extreme familiarity with honeypots before using them, more opensource honeypots should be built and deployed with advanced deception capabilities in their base configuration. This way, a wider range of individuals and organisations would be able to protect their networks, or study new attack methods being leveraged by hackers across the globe - without getting detected themselves."

In order to study this solution's feasibility, the creation of a lowinteraction honeypot has been carried out for network monitoring.

# 3. Materials & Method

#### 3.1. Architecture

A basic low-interaction honeypot has been created, with support for Hypertext Transfer Protocol (HTTP), Hypertext Transfer Protocol Secure (HTTPS), Secure Shell (SSH) and File Transfer Protocol (FTP) protocols. It is capable of logging all network traffic on its interfaces, parsing them, and sending summarised notifications on Slack Messenger - a messaging application built for and used extensively by businesses. The honeypot is capable of responding to attacker vulnerability probes and appears open to SSH connections, enabling the collection of login credentials being used from the attacker's side, for further analysis. As explained in Fig. (1), Python has been used as the programming language to deploy this honeypot on a virtual machine configured as a CentOS 8 x64 server, for minimal manual intervention over a period of multiple weeks of log collection. The honeypot system makes use of network monitoring tools on the server for the collection of the above-mentioned logs.

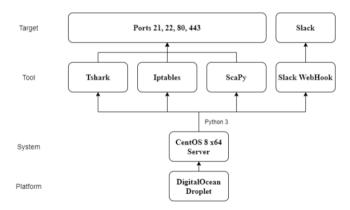


Fig. (1). Summary of Targets, Tools, the System and the Platform in use

# 3.2. Methodology

The research technique used for this study involved carrying out a comprehensive review of literature on honeypots. This required

gathering qualitative and quantitative data from a variety of sources - including books, journal papers, conference proceedings, and the Internet. Keywords such as "honeypot", "SSH logging", "network security", and "deception technology" were used for the same.

Parameters such as honeypot detectability, type (research/production), interaction (low/high), scalability (low cost/high cost), and implementation (software/hardware) were evaluated. After gathering this information, the sources were examined to see if they were pertinent, and duplicate information was eliminated. It was found that several sources featured more than one theme while the data was being gathered. In these situations, the prevailing subject matter was regarded as the principal theme of that source.

Finally, the advantages and disadvantages of each existing/proposed honeypot model were compared and combined to create a user-friendly, low-interaction honeypot that addresses

- support for detection of multiple communication protocols
- support for logging SSH credentials used via communicating with the system
- support for providing notifications of event summary via business channels

as discussed in this paper.

#### 3.2.1. Protocol Support Module

In order to capture all TCP network traffic at the default interface, Tshark - a network protocol analyzer - has been employed for FTP, SSH, HTTP and HTTPS logging on ports 21, 22, 80 and 443. Scapy - a packet manipulation program - has been used to check for FTP, SSH, HTTP and HTTPS SYN (synchronize) requests from any source and log each request with the source IP address, source port and destination port. Additionally, it replies with custom SYN-ACK (synchronize-acknowledge) packets to these requests - thus appearing vulnerable to insecure connections from attackers.

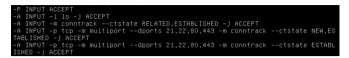


Fig. (2). Iptables INPUT chain rules instructing the system to enable communication on ports 21, 22, 80 and 443



Fig. (3). Iptables OUTPUT chain rules instructing the system to enable communication on ports 21, 22, 80 and 443

These packets are created on the basis of certain firewall rules, as seen in Fig. (2) and Fig. (3). If TCP packets from any source port on the outgoing interface have the RST (reset) flag set, the packets are dropped as RST indicates the need for connection termination. The RST iptables rule is dropped when the script stops running.

#### 3.2.2. SSH Credential Logging Module

By default, the SSH protocol logs SSH login attempts, regardless of whether or not authentication is successful. However, since it uses an encrypted tunnel for all communication, it isn't possible to read the data being sent and the local logs do not record the passwords being used. Therefore, it isn't possible to log the login credentials being used via SSH with its default configuration. In order to overcome this, the SSH configuration present on the server has been altered as required.

The altered configuration has been achieved by executing the following as the root user:

- 1) Uninstall the SSH server and download from the source.
- 2) Insert a logit() function in the SSH authentication file "authpasswd.c" at the location highlighted in Fig. (4).
- 3) Configure and install the SSH server as required.

```
    int
    sys_auth_passwd(struct ssh *ssh, const char *password)

[3]
[4]
[5]
[6]
[7]
[8]
         Authotxt *authotxt = ssh->authotxt;
# ---- INSERT HERE ----
logit("Honeypot_log: Username and Password -> %s,%s",authotxt->user,password);
# ^ INSERT HERE ^
          # ^ INSERT HERE auth_session_t *as;
static int expire_checked = 0;
[9]
[10
[11
                as = auth_usercheck(authctxt->pw->pw_name, authctxt->style, "auth-ssh",
  (char *)password);
[12]
[13]
[14]
[15]
[16]
[17]
[18]
[19]
                return (auth_close(as));
[20]
[21]
[22]
[23]
[24]
          ;
#elif !defined(CUSTOM SYS AUTH PASSWD)
          int
           sys_auth_passwd(struct ssh *ssh, const char *password)
                Authctxt *authctxt = ssh->authctxt;
[25]
[26]
[27]
                 # ---- INSERT HERE ----
logit("Honeypor_Logi Username and Fassword -> %s,%s",authotxt->user,password);
# ^ INSERT HERE ^
                # INSERT HERE "
struct passwd *pw = authctxt->pw;
char *encrypted_password, *salt = NULL;
28
[29]
[30]
[31]
321
               return encrypted_password != NULL &&
1331
[34]
[35]
[36]
                         trcmp(encrypted password, pw password) == 0;
          ,
#endif
```

```
Fig. (4). Credential logging function required in SSH server's password authentication file 'auth-passwd.c'
```

# 3.2.3. Notification Module

The need for timely, concise and easily accessible updates about possible attackers is extremely important for any organisation hosting a honeypot. Without it, there would be complete reliability in manually collecting traffic logs to detect and calculate all attempted connections to the honeypot. This would be slow, and prone to human errors. To accommodate this requirement, a Slack notification module has been included in this honeypot system. Slack is a messaging application used for team communications by businesses. It handles messages, files, third-party integrations such as Twitter, Dropbox, Google Docs, Trello, GitHub and dozens of other services all in one place. From large companies such as Pinterest, Airbnb and Shopify to smaller startups - all types of businesses use Slack - making it the ideal choice for an attack notification centre.

Slack's incoming webhook feature - a simple way to post messages from Slack applications to any channel - has been used to send updates about the number of connections attempted, to a Slack channel being used by the administrator (organisation). This has been achieved by reading all the source IP addresses from the traffic logs gathered by the honeypot, counting unique IP addresses found in the logs, and calculating which ones attempted the 110

maximum number of connections. Using the source IP addresses and the number of times they sent connection requests (Top 1, Top 2 or Top 3), messages are created and sent to the Slack channel.

#### 4. Observations

#### 4.1. Results

The analysis of the gathered network traffic logs reveals information such as the attackers' geographical location, protocols being used, timestamps of the attacks, etc. The success of this study has been determined by the running of the honeypot, the level of deception it provides, and the variety of data it successfully collects. These results aim to encourage developers to work on security solutions for all types and sizes of organisations, supporting future research that would provide insights into the current state of available solutions.

#### 4.1.1. Traffic Logging

The honeypot was deployed for 240 hours, from 21 October 2021 to 31 October 2021. Using the logs collected during this period, the following information was gathered:

28482, "Oct 21, 2021 29:17:33 SS8725380 UTC", 384.16.382.15,ette, TC", "443 + 38938 [SVN, ACX] Sep-0 Ack-1 Min+855335 Len-0 MS-51480 SAX FEM-1 MS-51024", United States
26031 "Oct 21, 2021 20 17 33 50070002 070" 44 277 100 59 410 102 30098 + 441 [AX1 Sent] 4 kt Windd/56 Land, United States
28484. "Oct 21, 2021 28:17:35.579500000 U/C".64.227.100.59.et50.15st.Client mallo.United Status
28405."Oct 21, 2021 28:17.33.580509725 07C", 104.16.102.15.ett0.TCP.443 + 38999 [ACK] Sept1 Ack-146 Min-67564 Lam-B.United States
28408."Oct 21, 2021 20-17-33-581955082 UTC",104.10.102.15_eth0.105v1.2, "Server Hello, Certificate Status, Server Key Sychampe, Server Hello Dove". United Status
28487,"0ct 21, 2011 20:17.33.581628330 UTC".04.227.186.59,ett0.TCP_38998 = 643 [ACK] Segr246 AcK+2866 Min+65616 Lenvil,United States
28688, "Oct 21, 2011 20:17:33.585683276 UTC", 64.227.186.59,eth0, TLSv1.2, "Client Key Excharge, Charge Cipher Spec, Encrypted Handshake Hessage", United States
28499, "Oct 21, 2821 28 17,33,588656897 UTC", 104,16,182,15,ettm, TCP.481 + 38998 (ACK) Securate Ackyldd Minus7588 (etu), United States
28898, "Oct 21, 2821 29-17-13, 585843161 UTC", 104, 16, 182, 15, ettm, TiSel J, "New Section Ticket, Change Claher Sect, Encrysted Handshake Rectage". United States
28491,"Oct 21, 2021 20:17:33.587145030 UTC".64,227.386.59.eth0.TL5v1.2.4pplication Data.United States
28492, "Oct 21, 2021 20:17-33.588063398 UTC", 504, 16.392.15.etb0.TCP.463 + 38998 [ACK] Seg-3096 Ack-613 Wirwidd000 [en-0.United States
28493, "Oct 21, 2021 20:17:33.640619733 UTC",104.16.102.15.et90.TLSv1.3.Application Outs.United Status
28494,"Oct 21, 2021 20:17:33.686945625 UTC".64.227.186.59.etb0.TCP.38998 + 443 [ACK] Secr613 Ack=3918 Win+04128 LenveLUnited States
28495, "Oct 21, 2021 20-17-33,722964812 UTC",64,227,180,59,eth0,TCP,"38996 = 643 [FIN, ACK] Seq-613 Ack-3918 Win-64128 Lenve", United Status
28496,"Oct 21, 2021 20:17:33.724142329 UTC",D04.16.182.15,eth0,TCP,"643 + 38998 (FIN, ACK) Seq=3918 Ark+614 Min+68608 Lev+0",United Status
28497,"Oct 21, 2021 20:17:33.724184667 UTC",64.227.186.50,eth0,TCP,38998 + 443 [ACK] Seq+614 Ack-3919 Win+64128 Lenvel,United States
28098,"Oct 21, 2021 20:18:20.302623562 UTC",108.08.09.124,eth0,TCP,01286 + 80 [SYN] Seque Wine5555 Level MSS-536,United States
28899,"Oct 21, 2021 20:18:29.302686332 UTC",64.227.506.59,ett0,TCP,"80 + 63206 [RST, dCK] Sepi1 Acks1 Winn0 Lenno",United States
28500,"Oct 21, 2021 20:10:25.392007501 UTC",64.227.106.59,eth0,TCP,"[TCP Port numbers reused] 00 + 43206 [SYN, ACK] Sequi Ack=1 Min#8192 Lenv0",United States
28591,"Oct 21, 2021 20:18:29.395897212 UTC",64.227.186.59,et00,TCP,"[TCP Retransmission] [TCP Port numbers reused] 80 + 43286 [SYN, 40X] Seoul 4ck+1 Min-8192 Len+0",UNLted States
28582,"0c1 21, 2021 20:18:29.605422914 UTC",198.98.49.124,e190,TCP,43126 - 80 [RS7] Sea-1 Win-0 Len-0,United States
28583, "Oct 21, 2021 28:18:29.600039451 UTC",198.96.49.124,e198, TC*,41286 + 10 [951] Sapt Mined Laned,Usited States
28584,"0C1 21, 2821 28:20:32.337581386.07C",95.217.31.46,4518,TCP,"445 + 9218 [596, 463] Sapet Ack+1 Mix+15384 Larvet 955+1468",doraina
28585,"Oct 21, 2021 29:21:21.004866885 UTC",171.251.20.236,eth0,TCF,51270 + 22 [5Y6] Seq=0 Micro64246 Lenv0 MS-12440 SetX,FEMP1 T5val-1982852539 T5ecr=0 MS-128,V1et Nam
20506,"Nrt 21, NV1 20-21-21.004091407 UTC",64.227.105.59,ett0,TO","12 + 51270 [5%, AK] Sept Acto1 MinetS100 Level MS-1408 SACK PERMIT Toul-110855330 M5-120",001ed States
28007, "Oct 21, 2001 20:71-23, 40164844 UTC", 44, 227, 488, 50, etch, [UD, "TCP Previous segment not ceptioned] [TCP Previous segment not 22 + 51278 [SN9, AX3 Seq1331031612 Abc1 MiniB192 [snn9", doi:14d States 28008, "Oct 23, 2001 20:71-23, 41654801 UTC", 44, 227, 488, 50, etch, [UD, "TTPP Previous segment not 24 + 13278 [SN9, AX3 Seq1331031612 Abc1, MiniB192 [snn9", doi:14d States 28008, "Oct 23, 2001 20:71-23, 41654801 UTC", 44, 327, 488, 50, etch, [UD, "TTPP Previous segment not 24 + 13278 [SN9, AX3] Seq1331031612 Abc1, MiniB192 [snn9", doi:14d States
20100 [CVC1 31, 2021 20:21(21, 619530571 UTC_66, 217, 100, 59, etc), (TCP Retranslation) [TCP Pert numbers reason] 22 + 51270 [UVB, AK] Sept3011051661 A(c)1 Knielly2 Lenter_initial States 20100 [CVC1 31, 2021 21, 217, 2164, 2164, 2164, 217, 2167, 217, 217, 217, 217, 217, 217, 217, 21
20100 YOL 23, 2001 BYLELEADORSDUE VIC, 121, 251, 26, 264, 269, 269, 269, 269, 269, 269, 269, 269
(2010) TXT 37, 2021 2011 (2.16796300) D1(7,17), 241-26, 26796 (SS) (11047) POTOS (SS) (2.167), V157 Base (2015) TXT 31, 2021 2017) (2.167974105 UTC -6.17) 105 (2.167074) (2.1670744) (2.1670744) (2.1670744) (
28513, "VCL 42, 2021 ab: 11:01.007/94176 UTC_A4.217.005.59; etcl.01.021 (State Laboration State Laboratio
2011, "UT 12, 2011 AV 1211, 00197001 UT _04.27.100.99," (UP 001-07-001) 27 = 31270 [PM, MA] Sup1 Advap Ni-05157 [sm-1] Test-19001002 [Sec-19002002] 27 = 31270 [PM, MA] Sup1 Advap Ni-05157 [sm-1] Test-19001002 [Sec-19002002] 27 = 31270 [PM, MA] Sup1 Advap Ni-05157 [sm-1] Test-19001002 [Sec-19002002] 27 = 31270 [PM, MA] Sup1 Advap Ni-05157 [sm-1] Test-19001002 [Sec-19002002] 27 = 31270 [PM, MA] Sup1 Advap Ni-05157 [sm-1] Test-19001002 [Sec-19002002] 27 = 31270 [PM, MA] Sup1 Advap Ni-05157 [sm-1] Test-19001002 [Sec-19002002] 27 = 31270 [PM, MA] Sup1 Advap Ni-05157 [sm-1] Test-19001002 [Sec-19002002] 27 = 31270 [Sm-1] Sup1 Advap Ni-05157 [sm-1] Test-19001002 [Sec-19002002] 27 = 31270 [Sm-1] Sup1 Advap Ni-05157 [sm-1] Test-19001002 [Sec-19002002] 27 = 31270 [Sm-1] Sup1 Advap Ni-05157 [sm-1] Test-19001002 [Sec-19002002] 27 = 31270 [Sm-1] Sup1 Advap Ni-05157 [sm-1] Test-19001002 [Sec-19002002] 27 = 31270 [Sm-1] Sup1 Advap Ni-05157 [sm-1] Test-19001002 [Sec-19002002] 27 = 31270 [Sm-1] Sup1 Advap Ni-05157 [sm-1] Test-19001002 [Sec-19002002] 27 = 31270 [Sm-1] Sup1 Advap Ni-05157 [sm-1] Test-19001002 [Sec-19002002] 27 = 31270 [Sm-1] Sup1 Advap Ni-05157 [sm-1] Test-19001002 [Sec-19002002] 27 = 31270 [Sm-1] Sup1 Advap Ni-05157 [sm-1] Test-19001002 [Sec-19002002] 27 = 31270 [Sm-1] Sup1 Advap Ni-05157 [sm-1] Test-19001002 [Sec-19002002] 27 = 31270 [Sm-1] Sup1 Advap Ni-05057 [sm-1] Sup1 Advap Ni-05057 [sm-1] Sup1 Advap Ni-05057 [sm-1] Sup1 Advap Ni-0507 [sm-1] Sup1
2011, VCT 11, 2011 AV12112.00790071 0(1,171.251.20.25009) (VT,1171 00) AAX 20009011 11179 422 (AAX) 500072 ACT1 BUTWARDS (Level 1044) 112171 11000411 11000110101 11000110101101
2015, VC 11, VX1 271112. AMMONG VC 101.121.2016.07.4000 (VC)1217 * 22 (AN) 30012 NOVACZ NEWSON (SWY 102120012047) SUPPLYS SUPPLYS AND
2010. No. 1. 2011 2012. The second state of
20137 Not 3. 2013 ADVID-112. TORONTO UNL AVAILABLE AND ADVID-10 FIELD ADVID-11 ADVID-112 ADVID-121 ADVI
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20104 Test 21, 2021 10 (112) 200000003 UTC 171 751 20 (20 atto: 5002 Client: Encrystal market (larette) Vist Rum
28525 "Oct 21, 2021 20:21:21.0006055666 U/C".54.227.100.59.etb0.107.22 + 51270 [ACK] Samp562 Ack+2413 Narrod418 [arrow [Sys]=15945335588 [Sacr-1982632623.3558] Sacr-1982632623.3558
28536. "Oct 21, 2021 20-21-21.8980773099 VIC" 171.253.30.256.eth0.55902.Client: Encrypted packet (len-64).Viet Num

Fig. (5). Logfile snippet

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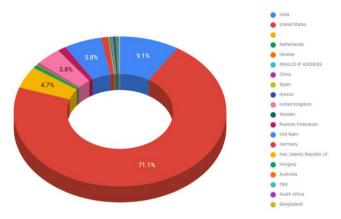


Fig. (6). Attack source (location) distribution

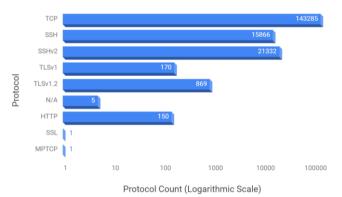


Fig. (7). Commonly exploited protocols

. Upon analysing the logs displayed in Fig. (5), it was observed that out of a total of 1,81,674 attempted connections, a strikingly large amount of traffic (71.1%) was generated from IP addresses mapped to the United States of America, while India reached the 9.1% mark - standing behind is Viet Nam at 5.8%. Other distinguishable locations included the United Kingdom (3.8%), the Russian Federation (1.2%) and the others (<=1%). Unidentifiable locations accounted for 4.7% of all traffic. While the difference in the amount of traffic generated by certain geographic locations may seem surprising in Fig. (6), factors such as technological advancement, infrastructure holding capacity and the usage of spoofed IP addresses or Virtual Private Networks must be kept in mind.

Overall, Fig. (7) shows a total of 1,43,285 TCP sessions, 2,13,323 SSHv2, and 15,866 SSH sessions. Across these sessions, the most commonly

exploited protocol was HTTPS, with 1040 unique requests. HTTP was used for 150 unique sessions, while other protocols were very rare.

#### 4.1.2. SSH Credential Logging

Following the custom SSH server configuration, the SSH local log file '/var/log/secure' not only contains records of attempted connections, but also the credentials used in those attempts - in plaintext, as evident in Fig. (8).

Oct	30	08:24:11	homeLab	sshd[6738]:	Ronewpot Log:	Username.	Password ->	root, sakuragi1
Oct	30	08:24:13	homeLab	sshd[6740]:	Honeypot_Log:	Username.	Password ->	root, roxy
Oct.	30	08:24:14	homeLab	sshd[67421:	Honeypot Log:	Username.	Password ->	root, romania
Oct	30	08:24:16	homeLab	sshd[67441:	Bonevpot Log	Username.	Password ->	root, rodel1
oct	30	08:24:17	homeLab	sshd[6747]:	Honeypot Log: Honeypot Log:	Username.	Password ->	root, robine
oct	30	08:24:19	homeLab	sshd[6749]:	Honeypot Log:	Username.	Password ->	root, ritale
Oct	30	08:24:20	homeLab	sshd[6751]:	Honevpot Log:	Username.	Password ->	root, riendutout
								root, rhbcnbyf
Oct	30	08:24:23	homeLab	sshd[6755]:	Honeypot Log:	Username,	Password ->	root, revolver
oct	30	08:24:25	homeLab	sshd[6757]:	Honeypot Log:	Username,	Password ->	root, revolver root, ramzi
oct		08:24:26	homeLab	sshd[6761]:	Honeypot Log:	Username,	Password ->	root, gwerty7
Oct		08:24:27	homeLab	sshd[6763]:	Honeypot Log:	Username,	Password ->	root, puntacana
Oct		09:36:59	homeLab	sshd[7255]:	Honeypot Log: Honeypot Log:	Username,	Password ->	devops, 1234
Oct		09:48:44	homeLab	sshd[7318]:	Boneypot Log:	Username,	Password ->	user, user
Oct		09:48:44	homeLab	sshd[7324]:	Honeypot Log:	Username,	Password ->	root, root
					Honeypot Log:			
Oct		09:48:45	homeLab	sshd[7322]:	Honeypot Log:	Username,	Password ->	admin, password
Oct		09:48:45	homeLab	sshd[7320]:	Honeypot Log: Honeypot Log:	Username,	Password ->	ubnt, ubnt
Oct		10:56:26	homeLab	sshd[8036]:		Username,	Password ->	root, root
Oct		10:56:26	homeLab	sshd[8035]:	Honeypot Log:	Username,	Password ->	vagrant, vagrant
Oct		10:56:26	homeLab	sshd[8038]:		Username,	Password ->	ubuntu, ubuntu
Oct		10:56:26	homeLab	sshd[8040]:	Honeypot Log:	Username,	Password ->	test, test
Oct	30	10:56:26	homeLab	sshd[8037]:		Username,	Password ->	oracle, oracle postgres, postgres
Oct	30	10:56:26	homeLab	sshd[8039]:	Honeypot_Log:	Username,	Password ->	postgres, postgres
								devops,123456
								devops, devops123
Oct	30	13:55:06	homeLab	sshd[8724]:	Honeypot Log:	Username,	Password ->	devops, password
Oct	30	15:06:52	homeLab	sshd[8919]:	Honeypot_Log:	Username,	Password ->	admin, admin admin, 010ctyQh243063uD
Oct	30	15:07:01	homeLab	sshd[8925]:	Honeypot_Log:	Username,	Password ->	admin,010ctyQh243063uD
Oct	30	15:07:12	homeLab	sshd[8931]:	Boneypot Log:	Username,	Password ->	system, OkwKcECs8qJP2Z
Oct	30	15:07:16	homeLab	sshd[8933]:	Honeypot Log:	Username,	Password ->	root, root
Oct	30	15:07:45	homeLab	sshd[8936]:	Boneypot_Log:	Username,	Password ->	user, user support, support admin, admin01
oct	30	15:07:46	homeLab	sshd[8939]:	Honeypot_Log:	Username,	Password ->	support, support
Oct	30	15:07:56	homeLab	sshd[8941]:	Honeypot Log: Honeypot Log:	Username,	Password ->	admin, admin01
Oct	30	15:08:31	nomeLab	ssna[8947]:	Honeypot Log:	Userhame,	Password ->	user1,123456
oct	30	15:21:53	nomeLab	sand[9042]:	Honeypot_Log	username,	Password ->	devops, devops@123 root, superman erp, erp
Oct	30	16:05:13	nomeLab	ssna[9128]:	Honeypot Log:	Userhame,	Password ->	root, superman
Oct	30	10:49:20	nomeLab	sshd[9245]:	Honeypot Log: Honeypot Log:	username,	Password ->	erp,erp
oct	30	19:04:10	nomeLab	sand[9686]:	Honeypot_Log	username,	Password ->	sky,MSIT#7pr0j21

Fig. (8). Filtered view of SSH server log file '/var/log/secure'

With over 315 unique usernames and 1233 unique passwords, the highest frequency was calculated for the credentials (in any combination) present in Table 2:

Usernames	Passwords
user	information
port	remote
root	admin
tracerlab	pi

ftp	oracle
rustserver	sync
webmaster	hyjx
mike	aaron
db2inst2	amy
install	m
reboot	support1
matt	Azureuser123
tmp	web
ems	carlos
dillon	Guest123
printer	bruce
ayden	xbian
belkinstyle	albaunio
paul	ts3
epg	alpha
pierre	nobody
ghost	youssra
new	vanesa
full	transformer
ts3	street1

guestuser	saluttoi
coach	romania
nobody	qwerty7
12qwaszx	superman

Table 2. Most frequently used usernames and passwords with SSH

#### 4.1.3. Slack Notifications

Useful in tracking down persistent attackers, the Slack notification module works to calculate the total number of connections attempted by IP addresses that interact with the honeypot frequently. Based on the logs collected during the abovementioned duration, the top 3 IP addresses that interacted with the honeypot made a total of 1,11,984 requests - as shown in Fig. (9), and the required information was sent as a message to the associated Slack channel.

> Maximum connections attempted by: 1. 64.227.186.59: 90841 connections 2. 117.236.129.78: 10939 connections 3. 173.255.236.224: 10204 connections

Fig. (9). Slack notification for the specified duration

#### 4.2. Analysis

Since SSH logs all attempted connections, the IP addresses associated with failed connections have also been recorded, along with the username. When required, this data may be analysed separately. Additionally, the SSH protocol allows authentication using keys, instead of passwords. Analysis of the log file shows that 53 unique IP addresses attempted key-based authentication a total of 2857 times, in addition to password-based authentication which has a total of 315 unique usernames with 1233 unique passwords in various combinations.

# 5. Conclusion

In this paper, we presented a user-friendly low-interaction honeypot. The honeypot is capable of running without manual intervention - once it has been deployed - and keeps track of each deployment session, without interfering with parallelly running processes (if any). The honeypot is capable of recording and analysing suspicious network traffic, as well as notifying the hosting organisation about the same. Additionally, it can collect login credentials used with SSH in plaintext, for a deeper insight into vulnerable keywords that may be blacklisted for increased security.

# **Challenges Faced**

- A large majority of currently available honeypots is built on outdated systems, with poor maintenance and irregular development cycles. They are predictable due to their limited adaptability and poor deception. Due to this, analysis of theory regarding fully functional honeypots that are user friendly enough to require minimal configuration, while being low interaction was difficult. However, by understanding the desirable aspects of multiple opensource honeypots, it was possible to integrate all the required functionality into one tool - while narrowing down on the exact architecture and tools needed for smooth functioning.
- 2) Default SSH logging of authentication attempts, while helpful, does not record passwords being used. Although this is a secure practice, it made the custom configuration of the SSH server on the honeypot a time-taking task. Taking inspiration from independent security researchers' attempts at implementing this idea [12], it was possible to create a solution that works with CentOS 8 x64 servers.

# **Future Scope**

In order to make the honeypot more comprehensive, support modules for analysing network requests captured with the traffic could be added. Doing so would allow researchers to get notified

about possible attack attempts such as HTTP-enabled backdoor installation. Additionally, platform support for a wider range of operating systems and environments could be added to reach a wider userbase.

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